

CHAPTER THREE

Section 3.1E

- 1.a. $2^1 = 2$
 c. $2^2 = 4$

2.a. \downarrow

E	$\sim \sim (E \ \& \ \sim E)$
T	F T T F F T
F	F T F F T F

c. \downarrow

A	J	$A \equiv [J \equiv (A \equiv J)]$
T	T	T T T T T T T
T	F	T T F T T F F
F	T	F T T F F F T
F	F	F T F F F T F

e. \downarrow

A	H	J	$[\sim A \vee (H \supset J)] \supset (A \vee J)$
T	T	T	F T T T T T T T T T
T	T	F	F T F T F F T T T F
T	F	T	F T T F T T T T T T
T	F	F	F T T F T F T T T F
F	T	T	T F T T T T T T F T T
F	T	F	T F T T F F F F F F F
F	F	T	T F T F T T T T F T T
F	F	F	T F T F T F F F F F F

g. \downarrow

A	B	$\sim (A \vee B) \supset (\sim A \vee \sim B)$
T	T	F T T T T F T F F T
T	F	F T T F T F T T T F
F	T	F F T T T T F T T F T
F	F	T F F F T T F T T F

i. \downarrow

B	E	H	$\sim (E \ \& \ [H \supset (B \ \& \ E)])$
T	T	T	F T T T T T T T T
T	T	F	F T T F T T T T T
T	F	T	T F F T F T T F F
T	F	F	T F F F T T T F F
F	T	T	T T F T F F F F T
F	T	F	F T T F T F F F T
F	F	T	T F F T F F F F F
F	F	F	T F F F T F F F F

k.

D	E	F	$\sim [D \ \& \ (E \ \vee \ F)]$	\equiv	$[\sim D \ \& \ (E \ \& \ F)]$
T	T	T	F	T	T
T	T	F	F	T	F
T	F	T	F	T	F
T	F	F	T	F	F
F	T	T	T	T	T
F	T	F	T	F	F
F	F	T	T	F	F
F	F	F	F	F	F

m.

A	H	J	$(A \ \vee \ (\sim A \ \& \ (H \ \supset \ J))) \ \supset \ (J \ \supset \ H)$
T	T	T	T
T	T	F	T
T	F	T	T
T	F	F	T
F	T	T	T
F	T	F	T
F	F	T	T
F	F	F	T

3.a.

A	B	C	$\sim [\sim A \ \vee \ (\sim C \ \vee \ \sim B)]$
F	T	T	F

c.

A	B	C	$(A \ \supset \ B) \ \vee \ (B \ \supset \ C)$
F	T	T	T

e.

A	B	C	$(A \ \equiv \ B) \ \vee \ (B \ \equiv \ C)$
F	T	T	F

g.

A	B	C	$\sim [B \ \supset \ (A \ \vee \ C)] \ \& \ \sim \sim B$
F	T	T	F

i.

A	B	C	$\sim [\sim (A \ \equiv \ \sim B) \ \equiv \ \sim A] \ \equiv \ (B \ \vee \ C)$
F	T	T	T

Section 3.2E

1.a. Truth-functionally indeterminate

$$\begin{array}{c|ccc}
 A & & \downarrow & \\
 & \sim A & \supset & A \\
 \hline
 \text{T} & \text{F} & \text{T} & \text{T} \\
 \text{F} & \text{T} & \text{F} & \text{F}
 \end{array}$$

c. Truth-functionally true

$$\begin{array}{c|cccccccc}
 A & & \downarrow & & & & & & \\
 & (A \equiv \sim A) & \supset & \sim(A \equiv \sim A) & & & & & \\
 \hline
 \text{T} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} \\
 \text{F} & \text{F} & \text{F} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T}
 \end{array}$$

e. Truth-functionally indeterminate

$$\begin{array}{cc|cccccccc}
 B & D & & \downarrow & & & & & & \\
 & & (\sim B \ \& \ \sim D) & \vee & \sim(B \ \vee \ D) & & & & \\
 \hline
 \text{T} & \text{T} & \text{F} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} & \text{T} \\
 \text{T} & \text{F} & \text{F} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} \\
 \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} \\
 \text{F} & \text{F} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{F} & \text{F}
 \end{array}$$

g. Truth-functionally indeterminate

$$\begin{array}{ccc|cccccccccccc}
 A & B & C & & \downarrow & & & & & & & & & \\
 & & & [(A \vee B) \ \& \ (A \vee C)] & \supset & \sim(B \ \& \ C) & & & & & & \\
 \hline
 \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} \\
 \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} \\
 \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} \\
 \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{F} \\
 \text{F} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} \\
 \text{F} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} \\
 \text{F} & \text{F} & \text{T} & \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} \\
 \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{F}
 \end{array}$$

i. Truth-functionally true

$$\begin{array}{cc|cccccccc}
 J & K & & \downarrow & & & & & & \\
 & & (J \vee \sim K) & \equiv & \sim\sim(K \supset J) & & & & & \\
 \hline
 \text{T} & \text{T} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{T} & \text{T} & \text{T} \\
 \text{T} & \text{F} & \text{T} & \text{T} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} & \text{T} & \text{T} \\
 \text{F} & \text{T} & \text{F} & \text{F} & \text{F} & \text{T} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} \\
 \text{F} & \text{F} & \text{F} & \text{F} & \text{T} & \text{T} & \text{F} & \text{F} & \text{T} & \text{F} & \text{F}
 \end{array}$$

e.

A	B	$(A \equiv B) \supset (A \supset B)$					
T	T	T	T	T	T	T	T
T	F	T	F	F	T	T	F
F	T	F	F	T	T	F	T
F	F	F	T	F	T	F	F

g.

A	B	C	$(A \supset B) \supset [(C \supset A) \supset (C \supset B)]$											
T	T	T	T	T	T	T	T	T	T	T	T	T	T	
T	T	F	T	T	T	T	F	T	T	T	F	T	T	
T	F	T	T	F	F	T	T	T	T	F	T	F	F	
T	F	F	T	F	F	T	F	T	T	T	F	T	F	
F	T	T	F	T	T	T	T	F	F	T	T	T	T	
F	T	F	F	T	T	T	F	T	F	T	F	T	T	
F	F	T	F	T	F	T	T	F	F	T	T	F	F	
F	F	F	F	T	F	T	F	T	F	T	F	T	F	

i.

A	B	$[(A \supset B) \& \sim B] \supset \sim A$							
T	T	T	T	T	F	F	T	T	F
T	F	T	F	F	F	T	F	T	F
F	T	F	T	T	F	F	T	T	F
F	F	F	T	F	T	T	F	T	F

k.

A	B	$A \supset [B \supset (A \supset B)]$					
T	T	T	T	T	T	T	
T	F	T	T	F	T	F	
F	T	F	T	T	F	T	
F	F	F	T	F	T	F	

m.

A	B	D	$(A \supset B) \supset [\sim B \supset \sim (A \& D)]$											
T	T	T	T	T	T	F	T	T	F	T	T	T		
T	T	F	T	T	T	T	F	T	T	T	T	F		
T	F	T	T	F	F	T	T	F	F	F	T	T		
T	F	F	T	F	F	T	T	F	T	T	T	F		
F	T	T	F	T	T	T	F	T	T	T	F	F		
F	T	F	F	T	T	T	F	T	T	T	F	F		
F	F	T	F	T	F	T	T	F	T	T	F	F		
F	F	F	F	T	F	T	T	F	T	T	F	F		

o.

A	B								
		↓							
~	(A ≡ B)	≡	(~	A	≡	B)		
T	T	F	T	T	T	T	F	T	F
T	F	T	T	F	F	T	F	T	T
F	T	T	F	F	T	T	T	F	T
F	F	F	F	T	F	T	T	F	F

4.a. Truth-functionally false

B	D						
		↓					
(B ≡ D)		&	(B	≡	~	D)	
T	T	T	T	T	F	T	F
T	F	T	F	F	F	T	T
F	T	F	F	T	F	F	T
F	F	F	T	F	F	F	T

c. Not truth-functionally false

A	B				
		↓			
A	≡	(B	≡	A)	
T	T	T	T	T	

e. Not truth-functionally false

C	D						
		↓					
[(C ∨ D) ≡ C]		⊃	~	C			
F	T	F	T	T	F	F	T

5.a. False. For example, while ‘(A ⊃ A)’ is truth-functionally true, ‘(A ⊃ A) & A’ is not.

c. True. There cannot be any truth-value assignment on which the antecedent is true and the consequent false because there is no truth-value assignment on which the consequent is false.

e. False. For example, although ‘(A & ~ A)’ is truth-functionally false, ‘C ∨ (A & ~ A)’ is not.

g. True. Since a sentence ~ **P** is false on a truth-value assignment if and only if **P** is true on the truth-value assignment, **P** is truth-functionally true if and only if ~ **P** is truth-functionally false.

i. False. For example, ‘(A ∨ ~ A)’ is truth-functionally true, but ‘(A ∨ ~ A) ⊃ B’ is truth-functionally indeterminate.

6.a. Yes On every truth-value assignment, **P** is true and **Q** is false. Hence **P** ≡ **Q** is false on every truth-value assignment. Therefore **P** ≡ **Q** is truth-functionally false.

c. No. Both 'A' and '~ A' are truth-functionally indeterminate, but 'A \vee ~ A' is truth-functionally true.

Section 3.3E

1.a. Not truth-functionally equivalent

A	B	\downarrow $\sim (A \ \& \ B)$	\downarrow $\sim (A \ \vee \ B)$
T	T	F	F
T	F	T	T
F	T	T	T
F	F	T	T

c. Truth-functionally equivalent

H	K	\downarrow $K \equiv H$	\downarrow $\sim K \equiv \sim H$
T	T	T	T
T	F	F	F
F	T	F	F
F	F	T	T

e. Truth-functionally equivalent

F	G	\downarrow $(G \supset F) \supset (F \supset G)$	\downarrow $(G \equiv F) \vee (\sim F \vee G)$
T	T	T	T
T	F	F	F
F	T	T	T
F	F	T	T

g. Not truth-functionally equivalent

H	J	K	\downarrow $\sim (H \ \& \ J) \equiv (J \equiv \sim K)$	\downarrow $(H \ \& \ J) \supset \sim K$
T	T	T	F	F
T	T	F	F	F
T	F	T	T	T
T	F	F	T	T
F	T	T	T	T
F	T	F	F	F
F	F	T	T	T
F	F	F	T	T

i. Not truth-functionally equivalent

A C D	$[A \vee \sim (D \& C)] \supset \sim D$	$[D \vee \sim (A \& C)] \supset \sim A$
T T T	T T F T T T F FT	T T F T T T F FT
T T F	T T T F F T T TF	F F F T T T T FT
T F T	T T T T F F F FT	T T T T F F F FT
T F F	T T T F F F T TF	F T T T F F F FT
F T T	F F F T T T T FT	T T T F F T T TF
F T F	F T T F F T T TF	F T T F F T T TF
F F T	F T T T F F F FT	T T T F F F T TF
F F F	F T T F F F T TF	F T T F F F T TF

k. Not truth-functionally equivalent

F G H	$F \vee \sim (G \vee \sim H)$	$(H \equiv \sim F) \vee G$
T T T	T T F T T FT	T F FT T T
T T F	T T F T T TF	F T FT T T
T F T	T T T F F FT	T F FT F F
T F F	T T F F T TF	F T FT T F
F T T	F F F T T FT	T T TF T T
F T F	F F F T T TF	F F TF T T
F F T	F T T F F FT	T T TF T F
F F F	F F F F T TF	F F TF F F

2.a. Truth-functionally equivalent

G H	$G \vee H$	$\sim G \supset H$
T T	T T T	FT T T
T F	T T F	FT T F
F T	F T T	TF T T
F F	F F F	TF F F

c. Truth-functionally equivalent

A D	$(D \equiv A) \& D$	$D \& A$
T T	T T T T T	T T T
T F	F F T F F	F F T
F T	T F F F T	T F F
F F	F T F F F	F F F

e. Not truth-functionally equivalent

A	$A \equiv (\sim A \equiv A)$	$\sim (A \supset \sim A)$
T	T F F T F T	T T F FT

3.a. Not truth-functionally equivalent

C: The sky clouds over.

N: The night will be clear.

M: The moon will shine brightly.

C	M	N	\downarrow C \vee (N & M)				\downarrow M \equiv (N & \sim C)				
T	T	T	T	T	T	T	T	F	T	F	FT
T	T	F	T	T	F	F	T	T	F	F	FT
T	F	T	T	T	T	F	F	F	T	T	FT
T	F	F	T	T	F	F	F	F	T	F	FT
F	T	T	F	T	T	T	T	T	T	T	TF
F	T	F	F	F	F	F	T	T	F	F	TF
F	F	T	F	F	T	F	F	F	F	T	TF
F	F	F	F	F	F	F	F	F	T	F	TF

c. Truth-functionally equivalent

D: The *Daily Herald* reports on our antics.

A: Our antics are effective.

A	D	\downarrow D \supset A			\downarrow \sim A \supset \sim D		
T	T	T	T	T	F	T	FT
T	F	F	T	T	F	T	TF
F	T	T	F	F	T	F	FT
F	F	F	T	F	T	F	TF

e. Not truth-functionally equivalent

M: Mary met Tom.

L: Mary liked Tom.

G: Mary asked George to the movies.

G	L	M	\downarrow (M & L) \supset \sim G				\downarrow (M & \sim L) \supset G					
T	T	T	T	T	T	F	FT	T	F	FT	T	T
T	T	F	F	F	T	T	FT	F	F	FT	T	T
T	F	T	T	F	F	T	FT	T	T	TF	T	T
T	F	F	F	F	F	T	FT	F	F	TF	T	T
F	T	T	T	T	T	T	TF	T	F	FT	T	F
F	T	F	F	F	T	T	TF	F	F	FT	T	F
F	F	T	T	F	F	T	TF	T	T	TF	F	F
F	F	F	F	F	F	T	TF	F	F	TF	T	F

4.a. Yes. **P** and **Q** have the same truth-value on every truth-value assignment. On every truth-value assignment on which they are both true, \sim **P** and

$\sim Q$ are both false, and on every truth-value assignment on which they are both false, $\sim P$ and $\sim Q$ are both true. It follows that $\sim P$ and $\sim Q$ are truth-functionally equivalent.

c. If P and Q are truth-functionally equivalent then they have the same truth-value on every truth-value assignment. On those assignments on which they are both true, the second disjunct of $\sim P \vee Q$ is true and so is the disjunction. On those assignments on which they are both false, the first disjunct of $\sim P \vee Q$ is true and so is the disjunction. So $\sim P \vee Q$ is true on every truth-value assignment.

5. No. Although the truth-tables for ' $A \vee B$ ' and ' $B \vee C$ ' have identical columns of truth-values, this does not show that the sentences are truth-functionally equivalent. To establish truth-functional equivalence we need to construct a single truth-table for both of the sentences:

A	B	C	↓ A ∨ B	B	↓ B ∨ C	C
T	T	T	T	T	T	T
T	T	F	T	T	T	F
T	F	T	T	T	T	T
T	F	F	T	T	F	F
F	T	T	F	T	T	T
F	T	F	F	T	T	F
F	F	T	F	F	T	T
F	F	F	F	F	F	F

As rows 4 and 7 show, ' $A \vee B$ ' and ' $B \vee C$ ' are not truth-functionally equivalent. Constructing separate truth-tables for these sentences rather than one table for both sentences does not show that these sentences have the same truth-values on *all* combinations of truth-values that their *combined* atomic components can have, and hence does not establish truth-functional equivalence.

6.a.

P	Q	↓ P ≡ Q	(P ⊃ Q)	↓ & (Q ⊃ P)	(P ⊃ Q) & (Q ⊃ P)
T	T	T	T	T	T
T	F	F	T	F	F
F	T	F	F	T	F
F	F	T	F	F	F

c.

P	Q	R	↓ P & (Q ∨ R)	(P & Q)	↓ ∨ (P & R)	(P & Q) ∨ (P & R)
T	T	T	T	T	T	T
T	T	F	T	T	T	T
T	F	T	T	F	T	T
T	F	F	T	F	F	F
F	T	T	F	F	T	F
F	T	F	F	F	F	F
F	F	T	F	F	T	F
F	F	F	F	F	F	F

e.

P	Q	\downarrow $\sim (P \equiv Q)$	\downarrow $(\sim P \equiv Q)$
T	T	F	T
T	F	T	F
F	T	T	T
F	F	F	F

Section 3.4E

1.a. Truth-functionally consistent

A	B	C	\downarrow $A \supset B$	\downarrow $B \supset C$	\downarrow $A \supset C$
T	T	T	T	T	T
T	T	F	T	F	F
T	F	T	F	T	T
T	F	F	F	F	F
F	T	T	F	T	T
F	T	F	F	F	F
F	F	T	F	T	T
F	F	F	F	F	F

c. Truth-functionally inconsistent

H	J	L	\downarrow $\sim [J \vee (H \supset L)]$	\downarrow $L \equiv (\sim J \vee \sim H)$	\downarrow $H \equiv (J \vee L)$
T	T	T	F	T	T
T	T	F	F	F	F
T	F	T	F	T	T
T	F	F	F	F	F
F	T	T	F	T	T
F	T	F	F	F	F
F	F	T	F	T	T
F	F	F	F	F	F

e. Truth-functionally inconsistent

H	J	\downarrow $(J \supset J) \supset H$	\downarrow $\sim J$	\downarrow $\sim H$
T	T	T	F	F
T	F	F	T	F
F	T	T	F	T
F	F	F	T	T

g. Truth-functionally consistent

A	B	C	↓	↓	↓
A	B	C	A	B	C
T	T	T	T	T	T
T	T	F	T	T	F
T	F	T	T	F	T
T	F	F	T	F	F
F	T	T	F	T	T
F	T	F	F	T	F
F	F	T	F	F	T
F	F	F	F	F	F

i. Truth-functionally consistent

A	B	C	↓	↓	↓
A	B	C	(A & B) ∨ (C ⊃ B)	~ A	~ B
T	T	T	T T T T T T T	FT	FT
T	T	F	T T T T F T T	FT	FT
T	F	T	T F F F T F F	FT	TF
T	F	F	T F F T F T F	FT	TF
F	T	T	F F T T T T T	TF	FT
F	T	F	F F T T F T T	TF	FT
F	F	T	F F F F T F F	TF	TF
F	F	F	F F F T F T F	TF	TF

2.a. Truth-functionally consistent

B	D	E	↓	↓
B	D	E	B ⊃ (D ⊃ E)	~ D & B
T	F	T	T T F T T	TF T T

c. Truth-functionally consistent

F	J	K	↓	↓
F	J	K	F ⊃ (J ∨ K)	F ≡ ~J
T	F	T	T T F T T	T T TF

e. Truth-functionally consistent

A	B	↓	↓
A	B	(A ⊃ B) ≡ (~ B ∨ B)	A
T	T	T T T T FT T T	T

3.a. Truth-functionally inconsistent

- S: Space is infinitely divisible.
 Z: Zeno's paradoxes are compelling.
 C: Zeno's paradoxes are convincing.

C	S	Z	S	⊃	Z	~	(C	∨	Z)	S
T	T	T	T	T	T	F	T	T	T	T
T	T	F	T	F	F	F	T	T	F	T
T	F	T	F	T	T	F	T	T	T	F
T	F	F	F	T	F	F	T	T	F	F
F	T	T	T	T	T	F	F	T	T	T
F	T	F	T	F	F	T	F	F	F	T
F	F	T	F	T	T	F	F	T	T	F
F	F	F	F	T	F	T	F	F	F	F

c. Truth-functionally consistent

- E: Eugene O'Neill was an alcoholic.
 P: Eugene O'Neill's plays show that he was an alcoholic.
 I: *The Iceman Cometh* must have been written by a teetotaler.
 F: Eugene O'Neill was a fake.

E	F	I	P	E	P	I	E	∨	F
T	T	T	T	T	T	T	T	T	T

e. Truth-functionally consistent

- R: The Red Sox will win next Sunday.
 J: Joan bet \$5.00 against the Red Sox.
 E: Joan will buy Ed a hamburger.

E	J	R	R	⊃	(J	⊃	E)	~	R	∧	~	E
F	T	F	F	T	T	F	F	T	F	T	T	F

4.a. First assume that $\{\mathbf{P}\}$ is truth-functionally inconsistent. Then, since \mathbf{P} is the only member of $\{\mathbf{P}\}$, there is no truth-value assignment on which \mathbf{P} is true; so \mathbf{P} is false on every truth-value assignment. But then $\sim \mathbf{P}$ is true on every truth-value assignment, and so $\sim \mathbf{P}$ is truth-functionally true.

Now assume that $\sim \mathbf{P}$ is truth-functionally true. Then $\sim \mathbf{P}$ is true on every truth-value assignment, and so \mathbf{P} is false on every truth-value assignment. But then there is no truth-value assignment on which \mathbf{P} , the only member of $\{\mathbf{P}\}$, is true, and so the set is truth-functionally inconsistent.

c. No. For example, 'A' and ' $\sim A$ ' are both truth-functionally indeterminate, but $\{A, \sim A\}$ is truth-functionally inconsistent.

Section 3.5E

1.a. Truth-functionally valid

A	H	J	$A \supset (H \ \& \ J)$	$J \equiv H$	$\sim J$	$\sim A$
T	T	T	T	T	F	F
T	T	F	F	F	T	F
T	F	T	F	F	F	F
T	F	F	F	F	T	F
F	T	T	T	T	F	T
F	T	F	F	F	T	T
F	F	T	F	F	F	T
F	F	F	F	F	T	T

c. Truth-functionally valid

A	D	G	$(D \equiv \sim G) \ \& \ G$	$(G \vee [(A \supset D) \ \& \ A]) \supset \sim D$	$G \supset \sim D$
T	T	T	F	T	F
T	T	F	T	T	F
T	F	T	F	F	T
T	F	F	T	F	T
F	T	T	F	T	F
F	T	F	T	F	F
F	F	T	T	F	T
F	F	F	F	F	T

e. Truth-functionally valid

C	D	E	$(C \supset D) \supset (D \supset E)$	D	C \supset E
T	T	T	T	T	T
T	T	F	F	T	F
T	F	T	F	F	T
T	F	F	F	F	F
F	T	T	T	T	F
F	T	F	F	T	F
F	F	T	F	F	F
F	F	F	T	F	F

g. Truth-functionally valid

G H	$(G \equiv H) \vee (\sim G \equiv H)$	$(\sim G \equiv \sim H) \vee \sim(G \equiv H)$
T T	T T T T F T F T	F T T F T T F T T T
T F	T F F T F T T F	F T F T F T T T T F F
F T	F F T T T F T T	T F F F T T T F F T
F F	F T F T T F F F	T F T T F T F F T F

i. Truth-functionally invalid

F G	$\sim \sim F \supset \sim \sim G$	$\sim G \supset \sim F$	$G \supset F$
T T	T F T T T F T	F T T F T	T T T
T F	T F T F F T F	T F F F T	F T T
F T	F T F T T F T	F T T T F	T F F
F F	F T F T F T F	T F T T F	F T F

2.a. Truth-functionally valid

J M	$(J \vee M) \supset \sim(J \& M)$	$M \equiv (M \supset J)$	$M \supset J$
T T	T T T F F T T T	T T T T T	T T T
T F	T T F T T T F F	F F F T T	F T T
F T	F T T T T F F T	T F T F F	T F F
F F	F F F T T F F F	F F F T F	F T F

c. Truth-functionally valid

A B	$A \supset \sim A$	$(B \supset A) \supset B$	$A \equiv \sim B$
T T	T F F T	T T T T T	T F F T
T F	T F F T	F T T F F	T T T F
F T	F T T F	T F F T T	F T F T
F F	F T T F	F T F F F	F F T F

e. Truth-functionally invalid

A B C	$A \& \sim[(B \& C) \equiv (C \supset A)]$	$B \supset \sim B$	$\sim C \supset C$
T F F	T T T F F F F F F T T	F T T F	T F F F

e. Truth-functionally valid

D: Computers can have desires.

E: Computers can have emotions.

T: Computers can think.

D	E	T	↓ T	≡	E	↓ E	⊃	D	↓ D	⊃	~	T	↓ ~	T
T	T	T	T	T	T	T	T	T	T	F	FT	F	T	
T	T	F	F	F	T	T	T	T	T	T	TF	T	F	
T	F	T	T	F	F	F	T	T	T	F	FT	F	T	
T	F	F	F	T	F	F	T	T	T	T	TF	T	F	
F	T	T	T	T	T	T	F	F	F	T	FT	F	T	
F	T	F	F	F	T	T	F	F	F	T	TF	F	T	
F	F	T	T	F	F	F	T	F	F	T	FT	F	T	
F	F	F	F	T	F	F	T	F	F	T	TF	F	T	

5. a. If $\{P\} \models Q$, then there is no truth-value assignment on which P is true and Q is false. If $\{Q\} \models P$, then there is no truth-value assignment on which Q is true and P is false. So there is no truth-value assignment on which P and Q have different truth-values, and therefore they are truth-functionally equivalent. Conversely, assume that P and Q are truth-functionally equivalent. Then Q is true on every truth-value assignment on which P is true, so $\{P\} \models Q$; and P is true on every truth-value assignment on which Q is true, so $\{Q\} \models P$.

c. Assume that $\{P\} \models Q$ and $\{Q\} \models R$. Then, by the first entailment, if P is true on a truth-value assignment, Q is also true on that assignment. By the second entailment, R must be true on that truth-value assignment as well. Therefore $\{P\} \models R$.

Section 3.6E

1.a. If $\{\sim P\}$ is truth-functionally inconsistent, then there is no truth-value assignment on which $\sim P$ is true (since $\sim P$ is the only member of its unit set). But then $\sim P$ is false on every truth-value assignment, so P is true on every truth-value assignment and is truth-functionally true.

c. If $\Gamma \cup \{\sim P\}$ is truth-functionally inconsistent, then there is no truth-value assignment on which every member of $\Gamma \cup \{\sim P\}$ is true. But $\sim P$ is true on a truth-value assignment if and only if P is false on that assignment. Hence

there is no truth-value assignment on which every member of Γ is true and \mathbf{P} is false. Hence $\Gamma \models \mathbf{P}$.

2.a. \mathbf{P} is truth-functionally true if and only if the set $\{\sim \mathbf{P}\}$ is truth-functionally inconsistent. But $\{\sim \mathbf{P}\}$ is the same set as $\emptyset \cup \{\sim \mathbf{P}\}$. So \mathbf{P} is truth-functionally true if and only if $\emptyset \cup \{\sim \mathbf{P}\}$ is truth-functionally inconsistent. But we have already seen, by previous results, that $\emptyset \cup \{\sim \mathbf{P}\}$ is truth-functionally inconsistent if and only if $\emptyset \models \mathbf{P}$. Hence \mathbf{P} is truth-functionally true if and only if $\emptyset \models \mathbf{P}$.

c. Assume that Γ is truth-functionally inconsistent. Then there is no truth-value assignment on which every member of Γ is true. Let \mathbf{P} be an *arbitrarily* selected sentence of SL . Then there is no truth-value assignment on which every member of Γ is true and \mathbf{P} false since there is no truth-value assignment on which every member of Γ is true. Hence $\Gamma \models \mathbf{P}$.

3.a. Since Γ is a truth-functionally consistent set there is at least one truth-value assignment on which every member of Γ is true. But \mathbf{P} is also true on such an assignment since a truth-functionally true sentence is true on every truth-value assignment. Hence on at least one truth-value assignment every member of $\Gamma \cup \{\mathbf{P}\}$ is true; so the set is truth-functionally consistent.

4.a. \mathbf{P} is either true or false on each truth-value assignment. On any assignment on which \mathbf{P} is true, \mathbf{Q} is true (because $\{\mathbf{P}\} \models \mathbf{Q}$) and so $\mathbf{Q} \vee \mathbf{R}$ is true. On any assignment on which \mathbf{P} is false, $\sim \mathbf{P}$ is true, \mathbf{R} is therefore also true (because $\{\sim \mathbf{P}\} \models \mathbf{R}$), and so $\mathbf{Q} \vee \mathbf{R}$ is true as well. Either way, then, $\mathbf{Q} \vee \mathbf{R}$ is true—so the sentence is truth-functionally true.

c. Assume that every member of $\Gamma \cup \Gamma'$ is true on some truth-value assignment. Then every member of Γ is true, and so \mathbf{P} is true (because $\Gamma \models \mathbf{P}$). Every member of Γ' is also true, and so \mathbf{Q} is true (because $\Gamma' \models \mathbf{Q}$). Therefore $\mathbf{P} \& \mathbf{Q}$ is true. So $\Gamma \cup \Gamma' \models \mathbf{P} \& \mathbf{Q}$.